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TITLE	PDP-8/E INSTRUCTION SIMULATORS FOR OTHER PDP-8S
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PDP-8/E Instruction Simulators for Other PDP-8 s

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Introduction

It used to be that PDP-8's came in four flavors: straight 8, 8/S, 8/I, and 8/L. Now there is a most marvellous fifth flavor: 8/E. The 8/E unlike the other four 8's has some extra instructions in the standard (non-EAE) set. This makes for easier programming, but messes up compatibility. Thus, although other 8 programs will run on an 8/E, 8/E programs may not run on other 8's.

There are three solutions to this problem:

- 1) don't use 8/E programs on your eight.
- 2) get rid of your other 8 and get an 8/E.
- 3) use some nice simulation subroutines like these.

The subroutines listed herein provide simulation for the functions of the PDP-8/E's BSW (byte swap) and standard MQ (non-EAE) microinstructions.

Advantages of Simulation Subroutines

- * They work.
- * They are called by a JMS I through a page zero pointer; thus the subroutines are callable from anywhere in 4K of core.
- * The subroutine calls are exactly as long as the microinstructions they replace: one location.
- * They are virtually independent of one another; thus they may be tucked into any page where there are a few words to spare, scattered all over core.
- * They're efficient (for simulation subroutines.)
- * They may be called by the same mnemonics as the actual instructions; no extensive program revision is necessary.

Disadvantages of Simulation Subroutines

- * They take up a little extra core.
- * They tie up page zero locations.
- * There is no provision for the use of multiple memory fields.
- * They're slower than the real thing. (If you're that anxious to save execution time, see option 2) above and ignore this stuff.)

Methods

I was inspired to write these routines by Robert H. Nagel (Ref. 1), who wrote a BSW simulator. ~~My~~ BSW simulator is an improved version of his original subroutine. The other microinstructions are simulated by maintaining a pseudo-MQ register on page zero. All the MQ operations are performed with respect to this pseudo-MQ.

Each simulator is invoked by using the standard microinstruction mnemonics BSW, MQA, MQL, CAM, ACL, SWP, and one I invented called ALC which is equivalent to CIA SWP which in turn equals CIA MQA MQL. These mnemonics are redefined to be JMS's to the appropriate simulator subroutines. Thus, assembling the subroutines with the user program causes the links to the simulators to be assembled into the program.

One location is required on page zero for the pseudo-MQ, plus one location for each subroutine to serve as a pointer for the JMS's.

The algorithms used for the simulations are relatively straightforward and are explained in comments in the listing. All algorithms are original with myself except for the exclusive OR algorithm (Ref. 2)

Restrictions

If the microinstructions are encountered in the forms MQA MQL, CLA MQL, etc. instead of SWP, CAM, etc, the instructions will not execute properly. This applies especially to QA SWP which has no single standard mnemonic. I have provided one: ALC. Any coding containing such combinations should be rewritten in terms of single mnemonics.

No provision is made for the possible desire to microprogram BSW with CLL, QML, STL, or IAC. The user may either write an extra simulator or simply use two separate instructions.

References

- 1) Robert H. Nagel, Letter to the Editor, DECUSCOPE, Vol. 10, No. 3, p. 23.
- 2) 1970 Small Computer Handbook, Digital Equipment Corporation, p. 28.


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/ PDP-3/E MICROINSTRUCTION SIMULATION ROUTINES
/
/ EACH SUBROUTINE IS INDEPENDENT OF THE OTHERS;
/ HOWEVER, NOTE THAT TO CONSERVE CORE LOCATIONS
/ XXCAM AND XXACL ARE USED AS TEMPORARY STORAGE
/ BY OTHER SUBROUTINES. IF XXCAM OR XXACL SHOULD
/ BE REMOVED A NEW LOCATION SHOULD BE DEFINED IF
/ NECESSARY FOR XXTMP1 OR XXTMP2.
/
/ NOTE THAT IF A SUBROUTINE IS REMOVED ITS PAGE
/ ZERO POINTER MAY ALSO BE REMOVED.
/
/ NOTE THAT THE MQL INSTRUCTION IS SIMULATED
/ SIMPLY BY A DCA RATHER THAN A JMS.
/
////////////////////////////////////
/ PAGE ZERO LOCATIONS
*170
MQL=DCA . / ANY PAGE ZERO LOCATIONS ARE OK
XXMQ, 0 / MQL (LOAD MQ FROM AC, CLEAR AC)
CAM=JMS I . / PSEUDO-MQ REGISTER
XXCAM / CAM (CLEAR AC AND MQ)
ACL=JMS I . / ACL (LOAD AC FROM MQ)
XXACL
SWP=JMS I . / SWP (SWAP AC AND MQ)
XXSWP
MQA=JMS I . / MQA (OR MQ INTO AC)
XXMQA
ALC=JMS I . / ALC (STANDS FOR AC LOAD
XXALC / AND CLEAR. THIS IS NOT A
/ STANDARD MNEMONIC AND SHOULD
/ BE DEFINED AS ALC=CLA SWP
/ FOR A PDP-3/E. LOAD AC FROM
/ MQ, CLEAR MQ)
/
BSW=JMS I . / BSW (BYTE SWAP. SWAP BITS 0-5
XXBSW / OF AC WITH BITS 6-11)
/
/ THE SUBROUTINES THEMSELVES MAY BE LOCATED ON ANY PAGE;
/ IF XXTMP1=XXACL AND XXTMP2=XXCAM THEN THE SUBROUTINES
/ SHOULD BE ARRANGED SO THAT THESE LOCATIONS CAN BE
/ ACCESSED, BUT OTHER THAN THIS THEY MAY BE TUCKED INTO
/ WHATEVER FREE AREAS EXIST IN CORE.
*600
/ CAM (CLEAR AC AND MQ)
/ ALGORITHM: QUITE STRAIGHTFORWARD
/ LENGTH: 4 LOCATIONS
/ EXECUTION TIME: 8 CYCLES
XXCAM, 0
CLA / CLEAR AC
DCA XXMQ / CLEAR MQ
JMP I XXCAM

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/ ACL (LOAD AC FROM MQ)
/ ALGORITHM: ABOUT AS STRAIGHTFORWARD AS FOR CAM
/ LENGTH: 4 LOCATIONS
/ EXECUTION TIME: 8 CYCLES
XXACL, Ø
    CLA                / CLEAR AC
    TAD XMQ           / GET MQ
    JMP I XXACL
/ SWP (SWAP AC AND MQ)
/ ALGORITHM: MOVE AC TO T1; MOVE MQ TO T2; MOVE
/      T1 TO MQ; MOVE T2 TO AC
/ LENGTH: 8 LOCATIONS
/ EXECUTION TIME: 17 CYCLES
XXSWP, Ø
    DCA XTMP1        / MOVE AC TO T1
    TAD XMQ
    DCA XTMP2        / MOVE MQ TO T2
    TAD XTMP1
    DCA XMQ          / MOVE T1 TO MQ
    TAD XTMP2        / MOVE T2 TO AC
    JMP I XXSWP
/ MQA (INCLUSIVE OR MQ INTO AC)
/ ALGORITHM: (AC IOR MQ) EQUALS (((NOT AC) AND MQ) + AC)
/      WHERE IOR, NOT, AND ARE LOGICAL OPERATORS. (SEE
/      1970 DEC SMALL COMPUTER HANDBOOK, PAGE 28.)
/ LENGTH: 7 LOCATIONS
/ EXECUTION TIME: 14 CYCLES
XXMQA, Ø
    DCA XTMP1        / AC
    TAD XTMP1
    CMA              / (NOT AC)
    AND XMQ          / ((NOT AC) AND MQ)
    TAD XTMP1        / (((NOT AC) AND MQ) + AC)
    JMP I XXMQA
/ ALC (AC LOAD AND CLEAR: LOAD AC FROM MQ, CLEAR MQ)
/ ALGORITHM: MOVE MQ TO T1, CLEAR MQ, MOVE T1 TO AC
/      (THIS SUBROUTINE COULD BE SHORTENED BY USING
/      XXSWP AT A COST OF EXECUTION TIME INCREASE.)
/ LENGTH: 7 LOCATIONS
/ EXECUTION TIME: 14 CYCLES
XXALC, Ø
    CLA
    TAD XMQ
    DCA XTMP1        / MOVE MQ TO T1
    DCA XMQ          / CLEAR MQ
    TAD XTMP1        / MOVE T1 TO AC
    JMP I XXALC

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/ BSW (BYTE SWAP: SWAP BITS 0-5 OF AC WITH BITS 6-11)
/ ALGORITHM: QUITE INGENIOUS (I'M PROUD OF THIS ONE.)
/   THIS ONE IS HARD TO DESCRIBE VERBALLY. THE
/   COMMENTS ON THE SUBROUTINE ILLUSTRATE WHERE
/   THE CONTENTS OF THE ORIGINAL AC AND LINK
/   APPEAR IN THE AC AND LINK AT EACH STEP.
/ LENGTH: 13 LOCATIONS
/ EXECUTION TIME: 21 CYCLES
/ NOTE THAT THE CONSTANT IS OFTEN USEFUL FOR OTHER
/ PARTS OF A USER PROGRAM.

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	LINK	AC
XXBSW, 0	/ L	XXX XXX YYY YYY
DCA XTMP1	/ L	000 000 000 000
RTR	/ 0	010 000 000 000
DCA XTMP2	/ 0	000 000 000 000
TAD XTMP1	/ 0	XXX XXX YYY YYY
AND XX7700	/ 0	XXX XXX 000 000
TAD XTMP1	/ X	XXX XX0 YYY YYY
RTL	/ X	XXX 0YY YYY YXX
RTL	/ X	X0Y YYY YYX XXX
TAD XTMP2	/ X	XLY YYY YYX XXX
RTL	/ L	YYY YYY XXX XXX
JMP I XXBSW	/	ALL DONE AT LAST!
XX7700, 7700	/	NIFTY CONSTANT

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/
/ THIS CONCLUDES THIS SET OF WONDERFUL PDP-8/E
/ INSTRUCTION SIMULATORS. I HOPE YOU HAVE
/ ENJOYED THEM.          - GUY L. STEELE JR.

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